

## Realizations and interconnections

- ▶ realizations and state-space notation
- ▶ poles, properness, SISO and MIMO
- ▶ equivalence of realizations
- ▶ operations: products, sums, inverses, feedback

# Realizations and interconnections

## State-space notation

Suppose

$$\hat{G}(z) = C(zI - A)^{-1}B + D$$

We define the notation

$$\left[ \begin{array}{c|c} A & B \\ \hline C & D \end{array} \right] = \hat{G}$$

This block-matrix notation means the *rational function*, and we can also write

$$\left[ \begin{array}{c|c} A & B \\ \hline C & D \end{array} \right] (z) = C(zI - A)^{-1}B + D$$

## Realization

given a rational transfer function  $G(s)$  (or  $G(z)$  in discrete time), a *realization* of  $G$  is a tuple  $(A, B, C, D)$  such that

$$G(s) = C(sI - A)^{-1}B + D$$

the *order* (or *dimension*) of the realization is  $n$ , the size of  $A$

- ▶ a realization is not unique: many different  $(A, B, C, D)$  can give the same  $G(s)$
- ▶ different realizations can have different dimensions
- ▶ *given*  $(A, B, C, D)$ , the transfer function is determined; going the other direction (from  $G$  to  $(A, B, C, D)$ ) is the *realization problem*

## Realization examples

**example 1:**  $G(s) = \frac{1}{s-1}$

verify:  $C(sI - A)^{-1}B + D = 1 \cdot \frac{1}{s-1} \cdot 1 + 0 = \frac{1}{s-1}$

realization:  $A = [1], B = [1], C = [1], D = 0$  (order 1)

**example 2:**  $G(s) = \frac{3}{s+2}$

realization:  $A = [-2], B = [1], C = [3], D = 0$  (order 1)

or:  $A = [-2], B = [3], C = [1], D = 0$  (also order 1)

both are valid realizations (the factorization of the residue  $3 = C \cdot B$  into  $C$  and  $B$  is not unique)

**example 3:**  $G(s) = 5$  (constant, no dynamics)

realization:  $D = [5]$  (order 0)

## Higher-order examples

**example 4:**  $G(s) = \frac{s+3}{(s+1)(s+2)} = \frac{s+3}{s^2+3s+2}$

one realization is

$$A = \begin{bmatrix} 0 & 1 \\ -2 & -3 \end{bmatrix}, \quad B = \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \quad C = [3 \quad 1], \quad D = 0$$

verify:  $C(sI - A)^{-1}B = [3 \quad 1] \frac{1}{s^2+3s+2} \begin{bmatrix} s+3 & 1 \\ -2 & s \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \frac{s+3}{s^2+3s+2}$

**example 5:** non-uniqueness of dimension

$G(s) = \frac{1}{s+1}$  has order-1 realization:  $A = [-1]$ ,  $B = [1]$ ,  $C = [1]$ ,  $D = 0$

but also an order-2 realization:  $A = \begin{bmatrix} -1 & 0 \\ 0 & -2 \end{bmatrix}$ ,  $B = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$ ,  $C = [1 \quad 0]$ ,  $D = 0$

(the mode at  $\lambda = -2$  does not appear in  $G(s)$ )

## MIMO example

**example 6:**  $A = \begin{bmatrix} -1 & 0 \\ 0 & -2 \end{bmatrix}$ ,  $B = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$ ,  $C = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}$ ,  $D = 0$

transfer function is  $2 \times 2$ :

$$G(s) = C(sI - A)^{-1}B = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} \frac{1}{s+1} & 0 \\ 0 & \frac{1}{s+2} \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} \frac{1}{s+1} & \frac{1}{s+2} \\ 0 & \frac{1}{s+2} \end{bmatrix}$$

- ▶  $G_{11}(s)$ : input 1 to output 1, pole at  $s = -1$
- ▶  $G_{12}(s)$ : input 2 to output 1, pole at  $s = -2$
- ▶  $G_{21}(s) = 0$ : input 1 does not affect output 2
- ▶  $G_{22}(s)$ : input 2 to output 2, pole at  $s = -2$

## Constructing realizations: partial fractions

**given:** SISO transfer function  $G(s) = \frac{n(s)}{d(s)}$ , strictly proper

expand in partial fractions (assuming distinct poles):

$$G(s) = \frac{r_1}{s - \lambda_1} + \frac{r_2}{s - \lambda_2} + \cdots + \frac{r_n}{s - \lambda_n}$$

read off the *diagonal realization*:

$$A = \begin{bmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \end{bmatrix}, \quad B = \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix}, \quad C = [r_1 \quad \cdots \quad r_n], \quad D = 0$$

verify:  $C(sI - A)^{-1}B = \sum_i r_i \cdot \frac{1}{s - \lambda_i} \cdot 1 = G(s)$

## Example: partial fraction realization

$$G(s) = \frac{s+3}{(s+1)(s+2)} = \frac{2}{s+1} + \frac{-1}{s+2}$$

diagonal realization:

$$A = \begin{bmatrix} -1 & 0 \\ 0 & -2 \end{bmatrix}, \quad B = \begin{bmatrix} 1 \\ 1 \end{bmatrix}, \quad C = [2 \quad -1], \quad D = 0$$

cf. the companion form realization from example 4:

$$A = \begin{bmatrix} 0 & 1 \\ -2 & -3 \end{bmatrix}, \quad B = \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \quad C = [3 \quad 1], \quad D = 0$$

these are two different realizations of the same  $G(s)$ ; they are related by a change of coordinates ( $T$  is the eigenvector matrix of the companion  $A$ )

## Constructing realizations: companion form

given: SISO  $G(s) = \frac{b_{n-1}s^{n-1} + \dots + b_0}{s^n + a_{n-1}s^{n-1} + \dots + a_0}$  (strictly proper)

the *companion form* realization is:

$$A = \begin{bmatrix} 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 1 & \cdots & 0 \\ \vdots & & & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & 1 \\ -a_0 & -a_1 & -a_2 & \cdots & -a_{n-1} \end{bmatrix}, \quad B = \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \\ 1 \end{bmatrix}$$

$$C = [b_0 \quad b_1 \quad \cdots \quad b_{n-1}], \quad D = 0$$

- ▶ denominator coefficients go in last row of  $A$
- ▶ numerator coefficients go in  $C$
- ▶  $\det(sI - A) = s^n + a_{n-1}s^{n-1} + \dots + a_0$  (characteristic polynomial read from  $A$ )

## Non-strictly-proper case

if  $G(s)$  is proper but not strictly proper:

$$G(s) = \frac{b_n s^n + b_{n-1} s^{n-1} + \dots + b_0}{s^n + a_{n-1} s^{n-1} + \dots + a_0}$$

perform polynomial long division:

$$G(s) = b_n + \frac{\check{b}_{n-1} s^{n-1} + \dots + \check{b}_0}{s^n + a_{n-1} s^{n-1} + \dots + a_0}$$

then  $D = b_n$  and apply companion form to the strictly proper remainder

**example:**  $G(s) = \frac{s+3}{s+1} = 1 + \frac{2}{s+1}$

realization:  $A = [-1]$ ,  $B = [1]$ ,  $C = [2]$ ,  $D = 1$

## Proper and strictly proper

$$G(s) = C(sI - A)^{-1}B + D$$

as  $s \rightarrow \infty$ :  $(sI - A)^{-1} \rightarrow 0$ , so  $G(s) \rightarrow D$

- ▶  $G$  is *proper* if  $G(\infty)$  is finite (always true for a state-space realization)
- ▶  $G$  is *strictly proper* if  $D = 0$ , i.e.,  $G(s) \rightarrow 0$  as  $s \rightarrow \infty$

strictly proper means no direct feedthrough from input to output

### SISO and MIMO:

- ▶  $m = p = 1$ : *single-input single-output* (SISO);  $G(s)$  is a scalar rational function
- ▶  $m > 1$  or  $p > 1$ : *multi-input multi-output* (MIMO);  $G(s) \in \mathbb{R}(s)^{p \times m}$  is a rational matrix

## Poles

the *poles* of  $G(s) = C(sI - A)^{-1}B + D$  are the values of  $s$  where  $G(s)$  is singular (*i.e.*, blows up)

since  $(sI - A)^{-1} = \text{adj}(sI - A) / \mathbf{det}(sI - A)$ , every pole is an eigenvalue of  $A$

- ▶ for SISO:  $G(s) = n(s)/d(s)$ ; poles are the roots of  $d(s)$  (after cancelling common factors with  $n(s)$ )
- ▶ some eigenvalues of  $A$  may cancel and not appear as poles of  $G(s)$  (we will see later that this happens exactly when the realization is not 'minimal')

## Changing coordinates

It is easy to see that if  $T$  is invertible, then

$$\left[ \begin{array}{c|c} A & B \\ \hline C & D \end{array} \right] = \left[ \begin{array}{c|c} T^{-1}AT & T^{-1}B \\ \hline CT & D \end{array} \right]$$

We call two realizations  $A_1, B_1, C_1, D_1$  and  $A_2, B_2, C_2, D_2$  *equivalent* if they have the same transfer function

$$\left[ \begin{array}{c|c} A_1 & B_1 \\ \hline C_1 & D_1 \end{array} \right] = \left[ \begin{array}{c|c} A_2 & B_2 \\ \hline C_2 & D_2 \end{array} \right]$$

equivalently (assuming  $D_1 = D_2$ ), any of:

- ▶ same impulse response:  $C_1 e^{A_1 t} B_1 = C_2 e^{A_2 t} B_2$  for all  $t \geq 0$
- ▶ same *Markov parameters*:  $C_1 A_1^k B_1 = C_2 A_2^k B_2$  for  $k = 0, 1, 2, \dots$

**note:** equivalent realizations can have different dimensions, different eigenvalues, and different state trajectories — they simply have the same input-output behavior

## Products

Suppose

$$G_1 = \left[ \begin{array}{c|c} A_1 & B_1 \\ \hline C_1 & D_1 \end{array} \right] \quad G_2 = \left[ \begin{array}{c|c} A_2 & B_2 \\ \hline C_2 & D_2 \end{array} \right]$$

Two realizations for the product are

$$\begin{aligned} G_1 G_2 &= \left[ \begin{array}{cc|c} A_1 & B_1 C_2 & B_1 D_2 \\ \hline 0 & A_2 & B_2 \\ \hline C_1 & D_1 C_2 & D_1 D_2 \end{array} \right] \\ &= \left[ \begin{array}{cc|c} A_2 & 0 & B_2 \\ \hline B_1 C_2 & A_1 & B_1 D_2 \\ \hline D_1 C_2 & C_1 & D_1 D_2 \end{array} \right] \end{aligned}$$

To prove this, use the fact that

$$\left[ \begin{array}{c|c} A & B \\ \hline 0 & C \end{array} \right]^{-1} = \left[ \begin{array}{c|c} A^{-1} & -A^{-1} B C^{-1} \\ \hline 0 & C^{-1} \end{array} \right]$$

## Sums

For the sum we have

$$G_1 + G_2 = \left[ \begin{array}{cc|c} A_1 & 0 & B_1 \\ 0 & A_2 & B_2 \\ \hline C_1 & C_2 & D_1 + D_2 \end{array} \right]$$

## Inverses

We can also invert a realization. If  $D$  is invertible, then

$$G^{-1} = \left[ \begin{array}{c|c} A - BD^{-1}C & -BD^{-1} \\ \hline D^{-1}C & D^{-1} \end{array} \right]$$

this means

$$\left[ \begin{array}{c|c} A - BD^{-1}C & -BD^{-1} \\ \hline D^{-1}C & D^{-1} \end{array} \right] \left[ \begin{array}{c|c} A & B \\ \hline C & D \end{array} \right] = I$$

where we multiply the two rational functions and cancel all common factors.

## Left Inverses

Similarly, if  $D$  is left invertible, then

$$\left[ \begin{array}{c|c} A - BD^\dagger C & -BD^\dagger \\ \hline D^\dagger C & D^\dagger \end{array} \right] \left[ \begin{array}{c|c} A & B \\ \hline C & D \end{array} \right] = I$$

## Stacking

$$\left[ \begin{array}{cc|cc} G_1 & G_2 & & \\ \hline A_1 & 0 & B_1 & 0 \\ 0 & A_2 & 0 & B_2 \\ \hline C_1 & C_2 & D_1 & D_2 \end{array} \right]$$

and

$$\left[ \begin{array}{c} G_1 \\ G_2 \end{array} \right] = \left[ \begin{array}{cc|cc} A_1 & 0 & B_1 & \\ \hline 0 & A_2 & B_2 & \\ C_1 & 0 & D_1 & \\ 0 & C_2 & D_2 & \end{array} \right]$$

## Feedback interconnection

plant  $G$  with controller  $K$  in negative feedback:  $u = r - Kw$ ,  $w = Gu$

the closed-loop from  $r$  to  $w$  is  $(I + GK)^{-1}G$

for  $G = \left[ \begin{array}{c|c} A & B \\ \hline C & 0 \end{array} \right]$  and static feedback  $K$ :

$$(I + GK)^{-1}G = \left[ \begin{array}{c|c} A - BKC & B \\ \hline C & 0 \end{array} \right]$$

for dynamic controller  $K = \left[ \begin{array}{c|c} A_k & B_k \\ \hline C_k & D_k \end{array} \right]$ :

$$(I + GK)^{-1}G = \left[ \begin{array}{cc|c} A - BD_kC & -BC_k & B \\ \hline B_kC & A_k & 0 \\ \hline C & 0 & 0 \end{array} \right]$$

(cf. separation principle: the augmented system with plant and observer-based controller)

**well-posedness:** requires  $I + D_kD$  invertible; automatic when either system is strictly proper

## Dimension of composed systems

every operation increases the state dimension (or leaves it unchanged):

$G_1 + G_2$                       dimension  $n_1 + n_2$

$G_1 G_2$                          dimension  $n_1 + n_2$

$[G_1 \quad G_2], \begin{bmatrix} G_1 \\ G_2 \end{bmatrix}$                       dimension  $n_1 + n_2$

$G^{-1}$                              dimension  $n$

feedback                         dimension  $n + n_k$

but these operations can result in realizations that are not the smallest possible

**example:**  $G_1 + G_2$  has dimension  $n_1 + n_2$ , but if  $G_1 = G_2$  then  $G_1 + G_2 = 2G_1$  has a simple realization with dimension  $n_1$